$S \rightarrow AB$ 

where A,B are in turn context free languages, is which construct?

- A union
- **B** intersection
- C concatenation
- D Kleene star
- E None of the above



 $S \rightarrow A \mid B$  where A,B are in turn context free languages, is which construct?

- A union
- **B** intersection
- C concatenation
- D Kleene star
- E None of the above



 $\Sigma$  is the alphabet

$$\textit{a} \in \Sigma$$

$$S \rightarrow As$$

$$As \rightarrow As \cdot a \mid a \mid \epsilon$$

where A,B are in turn context free languages, is which construct?

- A union
- **B** intersection
- C concatenation
- D Kleene star
- E None of the above



```
\Sigma is the alphabet
0 \in \Sigma
1 \in \Sigma
S \rightarrow 0S1 \mid \epsilon
is what language?
   A \{0^n1^n, n \in \mathbb{N}\}
   B \{0^i 1^j, i, j \in \mathbb{N}\}
   C \{0^i 1^j, i < j, i, j \in \mathbb{N}\}\
   D \{0^{j}1^{i}, i < j, i, j \in \mathbb{N}\}
    E None of the above
```



 $\Sigma$  is the alphabet  $0 \in \Sigma$   $1 \in \Sigma$   $S \to 0S0 \mid 1S1 \mid \epsilon$  is what language?

- A all palindromes
- B palindromes with an even number of symbols
- C palindromes with an odd number of symbols
- $D \{0^j 1^i, i < j, i, j \in \mathbb{N} \}$
- E None of the above



With the following hypothesis:

Context free grammars can construct the empty string, single tokens, unions, concatenations and Kleene stars.

Context free grammars can construct the language  $\{0^n1^n\}$  and  $\{ww^{\mathcal{R}}\}$ .

The class of context free grammars is closed under union, concatenation and Kleene star.

What can we conclude from this?

- A Some context free languages are infinite.
- B Some context free languages are regular languages.
- C Some context free languages are not regular languages.
- D B and C but not A.
- E A, B and C and not D.

Stop for a while



Can this rule be found in a context free grammar?

$$R_i \rightarrow A_i B_i i \geq 0$$

- A True
- B False



Can these rules be found in a context free grammar?

 $A \rightarrow BC$ 

 $B \rightarrow CA$ 

 $C \rightarrow AB$ 

A True

**B** False



Is this a context free grammar?

$$\Sigma = \{a, b, c\}$$

$$S \rightarrow \epsilon \mid A$$

$$A \rightarrow BC \mid a$$

$$B \rightarrow CA \mid b$$

$$C \rightarrow AB \mid c$$

- A True
- B False



Consider this grammar:

$$\Sigma = \{a, b\}$$

$$S \rightarrow AsAndBs$$

$$AsAndBs \rightarrow As \cdot AsAndBs \cdot Bs \mid a \mid b \mid \epsilon$$

$$As \rightarrow As \cdot a \mid a \mid \epsilon$$

$$Bs \rightarrow Bs \cdot b \mid b \mid \epsilon$$

Can you tell which sequence of substitutions generated aaaab?

- A Yes
- B No

stop here for a while



Let A and B be a context free languages, on  $\Sigma_A$  and  $\Sigma_B$ , respectively.

Let the start variable for a grammar for A be  $S_A$ , and the start variable for a grammar for B be  $S_B$ . Define

$$T = \{w_A w_b | w_A \in A \land w_B \in B\}.$$

Is T a context free language?

A Yes

B No



Can this grammar generate strings longer than any given positive integer, *m*?

$$\Sigma = \{a, b, c\}$$

$$S \rightarrow \epsilon \mid A$$

$$A \rightarrow BC \mid a$$

$$B \rightarrow CA \mid b$$

$$C \rightarrow AB \mid c$$

A Yes

B No



#### Consider this grammar:

$$\Sigma = \{a, b, c\}$$

$$S \rightarrow \epsilon \mid A$$

$$A \rightarrow BC \mid a$$

$$B \rightarrow CA \mid b$$

$$C \rightarrow AB \mid c$$

Suppose you could use each rule at most once. What is the longest string you could generate?

- A 1
- B 2
- **C** 3
- D > 3



If there is a length of string such that longer strings have to have been generated by using at least one rule more than once, are we going to get a pumping lemma from this?

A True

**B** False

stop for a while